Algorithm Engineering

Jens K. Mueller

jkm@informatik.uni-jena.de

Department of Mathematics and Computer Science Friedrich Schiller University Jena

Monday 2nd February, 2015

Theoretical Limits

Theoretical and Effective Performance

- Theoretical Performance
 - Number of nodes: 1
 - Number of CPUs: 1
 - Number of Cores: 1
 - CPU frequency: 2.66 GHz
 - Number of operations per cycle: 4 FLOP/cycle (single precision)

Theoretical (peak) performance

 $2.66\,\text{GHz}\cdot4\,\text{FLOP/cycle}\approx10.6\,\text{GFLOP/s}$

 Effective performance determines the number of operations per time

Theoretical and Effective Bandwidth

- Theoretical bandwidth
 - Memory clock: 1066 MHz
 - Type of memory: DDR2 (double data rate)
 - Number of channels: 2
 - Memory bus: 64 bit

 $1066 \text{ MHz} \cdot 2 \cdot 2 \cdot 64 \text{ bit}/8 = 34112 \text{ MB/s} \approx 34.1 \text{ GB/s}$

Effective bandwidth

$$(B_r+B_w)/t$$
,

where B_r is the number of bytes read, B_w number of bytes written and t the spent time.

System Information

- CPU
 - \$ lshw -C cpu
- Memory
 - \$ lshw -short -C memory

IO bound vs Compute bound

► Average #operations per in-/output

 $\frac{\text{number of operations}}{\text{size of in- and output}}$

 Number of operations needed to avoid CPU stalls (due to memory accesses)

Amdahl's Law

Fraction x of program optimized by a factor s, then



Assume T = 1 (unoptimized program's running time), then overall speedup is

$$\frac{T}{T_s} = \frac{1}{(1-x) + \frac{x}{s}}$$

- x = 1 yields overall speed up of s
- Strong scaling Overall speed up for fixed input size

Amdahl's Law (cont.)



Gustafson's Law

Embarrassingly parallel fraction x of program, then

$$T = (1 - x) + p \cdot x$$

Parallel running time is

$$T_p = (1-x) + x = 1$$

Hence, speed up is

$$\frac{T}{T_p} = (1-x) + p \cdot x = 1 + (p-1) \cdot x$$

- Parallel part grows linear with the number of processors
- Input sizes increase
- Weak scaling Overall speed up for increasing input size

